

CLOSURE SYSTEMS AND CLOSURE HYPERGROUPS

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SUNTO - Dato un sistema di chiusura (S, \mathfrak{C}) , sull'insieme S si può definire un'iperoperazione \cdot ponendo, per ogni $a, b \in S$, $a \cdot b = \langle a, b \rangle_{\mathfrak{C}}$, dove $\langle a, b \rangle_{\mathfrak{C}}$ è il minimo elemento di \mathfrak{C} a cui appartengono sia a che b . In questo lavoro noi studiamo questo tipo di iperoperazione, evidenziando diverse proprietà significative. Nei casi in cui l'iperoperazione in questione è associativa essa attribuisce ad S una struttura di ipergruppo, che noi chiamiamo *ipergruppo chiusura*.

ABSTRACT - If (S, \mathfrak{C}) is a closure system, then one can define on the set S a hyperoperation \cdot by setting, for any $a, b \in S$, $a \cdot b := \langle a, b \rangle_{\mathfrak{C}}$, where $\langle a, b \rangle_{\mathfrak{C}}$ is the minimum element of \mathfrak{C} containing a and b . In this paper we study such a type of hyperoperations and prove several interesting properties. Whenever the above hyperoperation is associative, then it gives S a hypergroup structure that we shall call *closure hypergroup*.

1. PRELIMINARIES AND RECALLS

A function $\cdot : S \times S \rightarrow \mathfrak{P}(S)$ is said a partial (binary) hyperoperation on S . If $x \cdot y \neq \emptyset$ ($x \cdot y := \cdot(x, y)$) for any $x, y \in S$, then one speaks of hyperoperation. For any $X, Y \subseteq S$ one can set $X \cdot Y := \bigcup_{x \in X, y \in Y} x \cdot y$ (hence $\emptyset \cdot Y = \emptyset = X \cdot \emptyset$). Thus one has also a binary operation on $\mathfrak{P}(S)$.

If $a \in S$ and $B \subseteq S$, then one usually writes respectively $a \cdot B$ and $B \cdot a$ instead of $\{a\} \cdot B$ and $B \cdot \{a\}$. It is obvious that $\bigcup_{x \in X} x \cdot Y = X \cdot Y = \bigcup_{y \in Y} X \cdot y$.

It is easy to verify that a partial hyperoperation on S is associative or commutative - with an obvious meaning of these terms - if and only if the corresponding operation on $\mathfrak{P}(S)$ is associative or commutative.

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Now we recall that a closure system on a set S is a subset \mathfrak{C} of the power set $\mathfrak{P}(S)$ which is closed under the arbitrary set intersection (in particular, $S = \bigcap \emptyset \in \mathfrak{C}$). One says also that (S, \mathfrak{C}) is a closure system.

For any $X \subseteq S$ one can consider the so called closure of X under \mathfrak{C} , given by the intersection of the elements of \mathfrak{C} including X , and represented by $\langle X \rangle_{\mathfrak{C}}$. If x_1, x_2, \dots, x_n are elements of S , then one writes $\langle x_1, x_2, \dots, x_n \rangle_{\mathfrak{C}}$ instead of $\langle \{x_1, x_2, \dots, x_n\} \rangle_{\mathfrak{C}}$ and says that $\langle x_1, x_2, \dots, x_n \rangle_{\mathfrak{C}}$ is finitely generated. The elements of \mathfrak{C} of the type $\langle x \rangle_{\mathfrak{C}}$ are said principal; moreover, if every element of \mathfrak{C} is principal, then (S, \mathfrak{C}) is said principal.

2. BINARY CLOSURE SYSTEMS

Through a closure system (S, \mathfrak{C}) one can define a commutative (binary) hyperoperation \cdot on S by setting, for any $a, b \in S$, $a \cdot b := \langle a, b \rangle_{\mathfrak{C}}$. We shall say that \cdot is a (binary) closure hyperoperation.

Remark 1. If \cdot is the above hyperoperation, since $a \cdot b = \langle a, b \rangle_{\mathfrak{C}}$ and $a \cdot b \subseteq \langle a \rangle_{\mathfrak{C}} \cdot b \subseteq \langle a \rangle_{\mathfrak{C}} \cdot \langle b \rangle_{\mathfrak{C}} \subseteq \langle a, b \rangle_{\mathfrak{C}}$ for any $a, b \in S$, then one has:

$$1) a \cdot b = \langle a \rangle_{\mathfrak{C}} \cdot b = a \cdot \langle b \rangle_{\mathfrak{C}} = \langle a \rangle_{\mathfrak{C}} \cdot \langle b \rangle_{\mathfrak{C}}.$$

Consequently, if $a \in S$ and $B \subseteq S$, then one has:

$$2) a \cdot B = \bigcup_{b \in B} a \cdot b = \bigcup_{b \in B} \langle a \rangle_{\mathfrak{C}} \cdot b = \langle a \rangle_{\mathfrak{C}} \cdot B (= B \cdot \langle a \rangle_{\mathfrak{C}}). \text{ In particular, if } \cdot \text{ is an associative hyperoperation and } x_1, x_2, \dots, x_n \in S, \text{ with } n > 2, \text{ then } x_1 \cdot x_2 \cdot \dots \cdot x_n = \langle x_1 \rangle_{\mathfrak{C}} \cdot \langle x_2 \rangle_{\mathfrak{C}} \cdot \dots \cdot \langle x_n \rangle_{\mathfrak{C}}.$$

Now, if \mathfrak{C} is a closure system on S , let \mathfrak{C}_2 be the set of the parts X of S which are closed under the hyperoperation \cdot associated to \mathfrak{C} (i.e.: $X \cdot X \subseteq X$). Thus we shall say that the set \mathfrak{C}_2 is a binary (or linear) closure system on S and call binary subspace (or pseudo-linear subspace) of (S, \mathfrak{C}) every element of \mathfrak{C}_2 . Obviously, \emptyset is a binary subspace of (S, \mathfrak{C}) ; moreover $\mathfrak{C} \subseteq \mathfrak{C}_2$.

Now we recall that a closure system (S, \mathfrak{C}) is said algebraic if, for any subset X of S and for any $x \in \langle X \rangle_{\mathfrak{C}}$, there is a finite subset F of X such that $x \in \langle F \rangle_{\mathfrak{C}}$. It is known that (S, \mathfrak{C}) is algebraic if and only if \mathfrak{C} is closed under the set union of the elements of any subset of \mathfrak{C} , which is upper directed (in particular, which is a chain) with respect to \subseteq . Therefore one can easily verify that \mathfrak{C}_2 is an algebraic closure system.

If $+$ is a partial hyperoperation on S , then it is clear that the set of the parts of S which are closed under $+$ is a binary closure system on S .

Remark 2. Let (S, \mathfrak{C}) and (S, \mathfrak{C}') be closure systems. If $\langle x, y \rangle_{\mathfrak{C}'} \subseteq \langle x, y \rangle_{\mathfrak{C}}$ for any $x, y \in S$, then $\mathfrak{C}_2 \subseteq \mathfrak{C}'_2$. Consequently, if $\mathfrak{C} \subseteq \mathfrak{C}'$, then $\mathfrak{C}_2 \subseteq \mathfrak{C}'_2$.

Remark 3. 1) If $x, y \in S$, then $\langle x, y \rangle_{\mathfrak{C}_2} \subseteq \langle x, y \rangle_{\mathfrak{C}}$. Moreover by definition of \mathfrak{C}_2 , since $x, y \in \langle x, y \rangle_{\mathfrak{C}_2}$, one has $\langle x, y \rangle_{\mathfrak{C}} \subseteq \langle x, y \rangle_{\mathfrak{C}_2}$. As a consequence one gets:

(a) $\langle x, y \rangle_{\mathfrak{C}} = \langle x, y \rangle_{\mathfrak{C}_2}$ (in particular, $\langle x \rangle_{\mathfrak{C}} = \langle x \rangle_{\mathfrak{C}_2}$). Therefore \mathfrak{C} and \mathfrak{C}_2 define the same binary closure hyperoperation, and hence $\mathfrak{C}_2 = (\mathfrak{C}_2)_2$.

(b) Since $\mathfrak{C}_2 = (\mathfrak{C}_2)_2$, \mathfrak{C} is a binary closure system if and only if $\mathfrak{C} = \mathfrak{C}_2$.

2) \mathfrak{C}_2 is the lower binary closure system on S including \mathfrak{C} . Indeed, if \mathfrak{C}' is a binary closure system and $\mathfrak{C} \subseteq \mathfrak{C}'$, then $\mathfrak{C}_2 \subseteq \mathfrak{C}'_2 = \mathfrak{C}'$.

3) If X is an element of \mathfrak{C}_2 , then $X \subseteq X \cdot X$, and hence $X = X \cdot X$.

Theorem 4. If (S, \mathfrak{C}) is a binary closure system, let C be a fixed element of \mathfrak{C} , and let $\mathfrak{C}' = \{Y \in \mathfrak{C} \mid Y = \emptyset \text{ or } C \subseteq Y\}$. Then (S, \mathfrak{C}') is a binary closure system.

Proof. It is obvious that (S, \mathfrak{C}') is a closure system. Thus, by Remark 3 (see (b) of property 1)), it is sufficient to prove that $\mathfrak{C}'_2 \subseteq \mathfrak{C}'$.

To this end, let X be a non empty element of \mathfrak{C}'_2 . Thus, for every $x, y \in X$, one has $C \subseteq \langle x, y \rangle_{\mathfrak{C}'} \subseteq X$ and $\langle x, y \rangle_{\mathfrak{C}} \subseteq \langle x, y \rangle_{\mathfrak{C}'} \subseteq X$. Therefore $X \in \mathfrak{C}'$. ■

2. PARA-NORMAL CLOSURE SYSTEMS

Now, in order to extend some interesting properties of the normal subgroups of a group, in this paragraph let us assume that \mathfrak{N} is a binary closure system on S , $+$ is a hyperoperation on S and \mathfrak{C} is the closure system of the subsets of S closed under $+$. Thus we shall indicate respectively with $\cdot_{\mathfrak{N}}$ and $\cdot_{\mathfrak{C}}$ the closure hyperoperations associated to \mathfrak{N} and to \mathfrak{C} .

We shall say that (S, \mathfrak{N}) is *para-normal* with respect to $+$ whenever the following condition holds:

$$(\circ) \quad \forall x, y \in S: \langle x \rangle_{\mathfrak{N}} + \langle y \rangle_{\mathfrak{N}} = \langle x, y \rangle_{\mathfrak{N}}.$$

Hence, since if $x, y \in S$ we have $x+y \subseteq \langle x \rangle_{\mathfrak{N}} + \langle y \rangle_{\mathfrak{N}} = \langle x, y \rangle_{\mathfrak{N}} = x \cdot_{\mathfrak{N}} y$, then the following property holds:

$$(\circ\circ) \quad \forall x, y \in S: x+y \subseteq x \cdot_{\mathfrak{N}} y.$$

And now let us assume that the binary closure system (S, \mathfrak{N}) is para-normal with respect to the hyperoperation $+$. Then we have the following theorems.

Theorem 5. for any $A, B \in \mathfrak{N}$ one has:

$$(*) \quad A+B = A \cdot_{\mathfrak{C}} B = A \cdot_{\mathfrak{N}} B.$$

Furthermore \mathfrak{N} is included in \mathfrak{C} .

Proof. Preliminarily let us remark that $A+B = A \cdot_{\mathfrak{C}} B$. In fact the following equalities are trivial:

$$\begin{aligned} A+B &= \bigcup_{a \in A, b \in B} \langle a \rangle_{\mathfrak{N}} + \langle b \rangle_{\mathfrak{N}} = \bigcup_{a \in A, b \in B} \langle a, b \rangle_{\mathfrak{N}} = \\ &= \bigcup_{a \in A, b \in B} a \cdot_{\mathfrak{N}} b = A \cdot_{\mathfrak{N}} B. \end{aligned}$$

In particular, one has $A+A = A \cdot_{\mathfrak{N}} A = A$. Thus \mathfrak{N} is included in \mathfrak{C} and hence $A \cdot_{\mathfrak{C}} B \subseteq A \cdot_{\mathfrak{N}} B$. As a consequence - since it is obvious that $A+B \subseteq A \cdot_{\mathfrak{C}} B$ - we get $A+B = A \cdot_{\mathfrak{C}} B = A \cdot_{\mathfrak{N}} B$. ■

Theorem 6. Let $A \in \mathfrak{C}$ be a union of elements of \mathfrak{N} . Then $A \in \mathfrak{N}$ ¹.

Proof. It is sufficient to prove that $A \cdot_{\mathfrak{N}} A \subseteq A$. Indeed one has:

$$\begin{aligned} A \cdot_{\mathfrak{N}} A &\subseteq \bigcup_{a, a' \in A} \langle a \rangle_{\mathfrak{N}} \cdot_{\mathfrak{N}} \langle a' \rangle_{\mathfrak{N}} = \bigcup_{a, a' \in A} \langle a, a' \rangle_{\mathfrak{N}} = \\ &= \bigcup_{a, a' \in A} \langle a \rangle_{\mathfrak{N}} + \langle a' \rangle_{\mathfrak{N}} \subseteq A+A \subseteq A. \quad \blacksquare \end{aligned}$$

¹ See the case of a subgroup which is a union of normal subgroups.

Remark 7. Let us point out that if the binary closure system (S, \mathfrak{N}) is para-normal with respect to an associative hyperoperation $+$ then, as in the case of normal subgroups of a group, also $\cdot_{\mathfrak{N}}$ is associative. Indeed, by Remark 1 and by Theorem 5, for any $a, b, c \in S$ we have:

$$\begin{aligned} a \cdot_{\mathfrak{N}} (b \cdot_{\mathfrak{N}} c) &= \langle a \rangle_{\mathfrak{N}} \cdot_{\mathfrak{N}} (\langle b, c \rangle_{\mathfrak{N}}) = \langle a \rangle_{\mathfrak{N}} + (\langle b \rangle_{\mathfrak{N}} + \langle c \rangle_{\mathfrak{N}}) = \\ &= (\langle a \rangle_{\mathfrak{N}} + \langle b \rangle_{\mathfrak{N}}) + \langle c \rangle_{\mathfrak{N}} = \langle a, b \rangle_{\mathfrak{N}} \cdot_{\mathfrak{N}} \langle c \rangle_{\mathfrak{N}} = (a \cdot_{\mathfrak{N}} b) \cdot_{\mathfrak{N}} c. \end{aligned}$$

In the meantime $\cdot_{\mathfrak{C}}$ can be not associative, as in most groups $(S, +)$ in which $+$ is a non commutative operation.

3. ASSOCIATIVE CLOSURE SYSTEMS

Now let \cdot be the hyperoperation associated to a given closure system (S, \mathfrak{C}) ; hence, for any $x, y \in S$, x and y belong to the hyperproduct $x \cdot y$. As a consequence, if \cdot is associative, then \cdot gives S a structure of commutative hypergroup (in the sense of [1], p. 8). Therefore we shall say that (S, \mathfrak{C}) is *associative* and (S, \cdot) is a *closure hypergroup*.

Furthermore, we shall say that (S, \mathfrak{C}) is *3-strong associative* if, for any $x, y, z \in S$, $x \cdot (y \cdot z) = \langle x, y, z \rangle_{\mathfrak{C}} (= z \cdot (x \cdot y) = (x \cdot y) \cdot z)$. In such a case, since $x \cdot (y \cdot z) \subseteq \langle x, y, z \rangle_{\mathfrak{C}_2} \subseteq \langle x, y, z \rangle_{\mathfrak{C}}$, then one has $\langle x, y, z \rangle_{\mathfrak{C}} = \langle x, y, z \rangle_{\mathfrak{C}_2}$.

More generally, given a natural numbers $n \geq 2$, we shall say that (S, \mathfrak{C}) is *n-strong associative* if it is associative and $x_1 \cdot x_2 \cdot \dots \cdot x_n = \langle x_1, x_2, \dots, x_n \rangle_{\mathfrak{C}}$ for any $x_1, x_2, \dots, x_n \in S$. Furthermore, we shall say that (S, \mathfrak{C}) is *finitely strong associative* if it is n-strong associative for any natural numbers $n \geq 2$.

Obviously, a closure system (S, \mathfrak{C}) is 2-strong associative if and only if it is associative. Furthermore, if m and n are natural numbers such that $2 \leq m < n$ and (S, \mathfrak{C}) is n-strong associative, then (S, \mathfrak{C}) is m-strong associative. In fact, if one set $x_m = x_{m+1} = \dots = x_n$, then (see 1) of Remark 1) one gets ²:

$$x_1 \cdot x_2 \cdot \dots \cdot x_{m-1} \cdot x_m = x_1 \cdot x_2 \cdot \dots \cdot x_{m-1} \cdot \langle x_m \rangle_{\mathfrak{C}} =$$

² We recall that $\langle x \rangle_{\mathfrak{C}} \cdot \langle x \rangle_{\mathfrak{C}} = \langle x \rangle_{\mathfrak{C}}$ for any $x \in S$.

$$\begin{aligned}
&= x_1 \cdot x_2 \cdot \dots \cdot x_{m-1} \cdot \langle x_m \rangle_{\mathfrak{C}} \cdot \langle x_{m+1} \rangle_{\mathfrak{C}} \cdot \dots \cdot \langle x_n \rangle_{\mathfrak{C}} = \\
&= x_1 \cdot x_2 \cdot \dots \cdot x_m \cdot x_{m+1} \cdot \dots \cdot x_n = \langle x_1, x_2, \dots, x_n \rangle_{\mathfrak{C}} = \\
&= \langle x_1, x_2, \dots, x_m \rangle_{\mathfrak{C}}.
\end{aligned}$$

Theorem 8. Let n be a natural number greater than 1. Then a closure system $(\mathbf{S}, \mathfrak{C})$ is n -strong associative if and only if the following property holds³:

$$(*) \quad \forall x_1, x_2, \dots, x_n \in \mathbf{S}: \langle x_1, x_2, \dots, x_n \rangle_{\mathfrak{C}} = x_1 \cdot \langle x_2, \dots, x_n \rangle_{\mathfrak{C}}.$$

Proof. The assertion is obvious if $n = 2$ or $n = 3$. Thus let $n > 3$. If $(\mathbf{S}, \mathfrak{C})$ is n -strong associative then, since it is also $(n-1)$ -strong associative, we have:

$$\langle x_1, x_2, \dots, x_n \rangle_{\mathfrak{C}} = x_1 \cdot x_2 \cdot \dots \cdot x_n = x_1 \cdot \langle x_2, \dots, x_n \rangle_{\mathfrak{C}}.$$

Conversely, let the condition $(*)$ hold. Then it holds also with n replaced by a natural number m such that $2 < m < n$. In fact we can set $x_m = x_{m+1} = \dots = x_n$. Thus, by setting $m = 3$, we have that $(\mathbf{S}, \mathfrak{C})$ is associative.

As an immediate consequence, by induction, we get $\langle x_1, x_2, \dots, x_n \rangle_{\mathfrak{C}} = x_1 \cdot x_2 \cdot \dots \cdot x_n$. ■

Theorem 9. Let $(\mathbf{S}, \mathfrak{C})$ be a closure system and let $a_1, a_2, \dots, a_n, b \in \mathbf{S}$. If $\langle a_2, \dots, a_n \rangle_{\mathfrak{C}} = \langle b \rangle_{\mathfrak{C}}$, then $a_1 \cdot \langle a_2, \dots, a_n \rangle_{\mathfrak{C}} = \langle a_1, a_2, \dots, a_n \rangle_{\mathfrak{C}}$.

Proof. Indeed (see 1) of Remark 1), $a_1 \cdot \langle a_2, \dots, a_n \rangle_{\mathfrak{C}} = a_1 \cdot \langle b \rangle_{\mathfrak{C}} = a_1 \cdot b = \langle a_1, b \rangle_{\mathfrak{C}} = \langle a_1, a_2, \dots, a_n \rangle_{\mathfrak{C}}$; whence the thesis. ■

Remark 10. If the finitely generated and non empty elements of a closure system \mathfrak{C} are principal, then (by Theorem 6) the condition $(*)$ of Theorem 8 is true for any natural number n , and hence $(\mathbf{S}, \mathfrak{C})$ is finitely strong associative. In particular, the closure system \mathfrak{I} of the ideals of a semilattice (\mathbf{S}, \cup) ⁴ is a finitely strong associative closure system. In fact one can immediately verify that, for any $x_1, x_2, \dots, x_n \in \mathbf{S}$, the ideal generated by x_1, x_2, \dots, x_n is equal to $\langle x_1 + x_2 + \dots + x_n \rangle_{\mathfrak{I}}$.

³ If $n = 2$ then, by 1) of Remark 1, property $(*)$ is true even if $(\mathbf{S}, \mathfrak{C})$ is non associative.

⁴ A semilattice is a structure (\mathbf{S}, \cup) , where \cup is an idempotent, commutative and associative binary operation; an ideal is a subset \mathbf{B} of \mathbf{S} closed under \cup such that, for any $x \in \mathbf{S}$ and $x' \in \mathbf{B}$, if $x \leq x'$ (i.e.: $x \cup x' = x'$), then $x \in \mathbf{B}$.

Theorem 11. Let (S, \mathfrak{C}) be an algebraic and associative closure system. Then \mathfrak{C} is finitely strong associative if and only if $\mathfrak{C} \cup \{\emptyset\} = \mathfrak{C}_2$.

Proof. Let (S, \mathfrak{C}) be finitely strong associative. Thus, since $\mathfrak{C} \cup \{\emptyset\} \subseteq \mathfrak{C}_2$, in order to prove that $\mathfrak{C} \cup \{\emptyset\} = \mathfrak{C}_2$ it is sufficient to verify that if X is a non empty element of \mathfrak{C}_2 , then $\langle X \rangle_{\mathfrak{C}} = X$ (hence X is also an element of \mathfrak{C}).

Since $X \subseteq \langle X \rangle_{\mathfrak{C}}$, we only have to verify that $\langle X \rangle_{\mathfrak{C}} \subseteq X$. Thus let $x' \in \langle X \rangle_{\mathfrak{C}}$ and (by the hypothesis that (S, \mathfrak{C}) is algebraic) let us consider $x_1, \dots, x_n \in X$ such that $x' \in \langle x_1, \dots, x_n \rangle_{\mathfrak{C}}$. As \mathfrak{C} is finitely strong associative, $\langle x_1, \dots, x_n \rangle_{\mathfrak{C}} = x_1 \cdot \dots \cdot x_n \subseteq X$, and hence $x' \in X$.

On the contrary, let $\mathfrak{C} \cup \{\emptyset\} = \mathfrak{C}_2$. Hence \mathfrak{C} and \mathfrak{C}_2 determine the same hyperoperation \cdot ; moreover $\langle X \rangle_{\mathfrak{C}_2} = \langle X \rangle_{\mathfrak{C}}$ for any non empty subset X of S , hence \mathfrak{C} is finitely strong associative if and only if \mathfrak{C}_2 is finitely strong associative. Thus let us verify that if $x_1, x_2, \dots, x_n \in S$, then $x_1 \cdot x_2 \cdot \dots \cdot x_n = \langle x_1, x_2, \dots, x_n \rangle_{\mathfrak{C}_2}$.

Indeed, since $x_1, x_2, \dots, x_n \in x_1 \cdot x_2 \cdot \dots \cdot x_n \subseteq \langle x_1, x_2, \dots, x_n \rangle_{\mathfrak{C}_2}$ and $\langle x_1, x_2, \dots, x_n \rangle_{\mathfrak{C}_2}$ is the minimum element of \mathfrak{C}_2 containing $\{x_1, x_2, \dots, x_n\}$, it is sufficient to point out that, by associativity and by commutativity, one has (cf. Remark 1):

$$\begin{aligned} (x_1 \cdot x_2 \cdot \dots \cdot x_n) \cdot (x_1 \cdot x_2 \cdot \dots \cdot x_n) &= x_1 \cdot x_1 \cdot x_2 \cdot x_2 \cdot \dots \cdot x_n \cdot x_n = \\ &= \langle x_1 \rangle_{\mathfrak{C}_2} \cdot \langle x_2 \rangle_{\mathfrak{C}_2} \cdot \dots \cdot \langle x_n \rangle_{\mathfrak{C}_2} = x_1 \cdot x_2 \cdot \dots \cdot x_n. \quad \blacksquare \end{aligned}$$

BIBLIOGRAPHY

1. P.G. Corsini, *Prolegomena of hypergroup theory*, Aviani, Tricesimo (UD; I), 1992.